

CSE 114A

Foundations of Programming Languages

Type classes

Overloading Operators: Arithmetic

The `+` operator works for a bunch of different types.

For `Integer`:

```
λ> 2 + 3
```

```
5
```

for `Double` precision floats:

```
λ> 2.9 + 3.5
```

```
6.4
```

Overloading Operators: Arithmetic

Similarly we can *compare* different types of values

```
λ> 2 == 3
```

```
False
```

```
λ> [2.9, 3.5] == [2.9, 3.5]
```

```
True
```

```
λ> ("cat", 10) < ("cat", 2)
```

```
False
```

```
λ> ("cat", 10) < ("cat", 20)
```

```
True
```

Ad-Hoc Overloading

Seems unremarkable?

Languages since the dawn of time have supported “operator overloading”

- To support this kind of **ad-hoc polymorphism**
- Ad-hoc: “created or done for a particular purpose as necessary.”

You really **need** to *add* and *compare* values of *multiple* types!

Haskell has no caste system

No distinction between operators and functions

- All are first class citizens!

But then, what type do we give to *functions* like `+` and `==` ?

Individual: Plus type

Which of the following would be appropriate types for (+) ?

(A) (+) :: Integer -> Integer -> Integer

(B) (+) :: Double -> Double -> Double

(C) (+) :: a -> a -> a

(D) *All* of the above

(E) *None* of the above



<https://tiny.cc/cse116-plus-type-ind>

Individual: Plus type

Which of the following would be appropriate types for (+) ?

(A) (+) :: Integer -> Integer -> Integer

(B) (+) :: Double -> Double -> Double

(C) (+) :: a -> a -> a

(D) *All* of the above

(E) *None* of the above



<http://tiny.cc/cse116-plus-type-grp>

Haskell has no caste system

`Integer -> Integer -> Integer` is bad because?

- Then we cannot add `Doubles`!

Haskell has no caste system

`Double -> Double -> Double` is bad because?

- Then we cannot add `Integer`!

Haskell has no caste system

$a \rightarrow a \rightarrow a$ is bad because?

- That doesn't make sense, e.g. to add two `Bool` or two `[Int]` or two functions!

Type Classes for Ad Hoc Polymorphism

Haskell solves this problem with an *insanely slick* mechanism called typeclasses, introduced by [Wadler and Blott](#)

How to make *ad-hoc* polymorphism less *ad hoc*

Philip Wadler and Stephen Blott
University of Glasgow*

October 1988

Qualified Types

To see the right type, lets ask:

```
λ> :type (+)
```

```
(+) :: (Num a) => a -> a -> a
```

We call the above a **qualified type**. Read it as +

- takes in two a values and returns an a value for any type a that

- *is a* Num or
- *implements* the Num interface or
- *is an instance of* a Num.

The name Num can be thought of as a *predicate* or *constraint* over types

Some types are Num

- Examples include **Integer**, **Double** etc
- Any such values of those types can be passed to **+**.

Other types are not Nums

Examples include `Char`, `String`, functions etc,

- Values of those types *cannot* be passed to `+`.

```
λ> True + False
```

```
<interactive>:15:6:
```

```
  No instance for (Num Bool) arising from a  
use of '+'
```

```
  In the expression: True + False
```

```
  In an equation for 'it': it = True + False
```

Type Class is a Set of Operations

A ***typeclass*** is a collection of operations (functions) that must exist for the underlying type.

The Eq Type Class

The simplest typeclass is perhaps, `Eq`

```
class Eq a where
  (==) :: a -> a -> Bool
  (/=) :: a -> a -> Bool
```

A type `a` is *an instance of* `Eq` if there are two functions

- `==` and `/=`

That determine if two `a` values are respectively *equal* or *unequal*.

The Show Type Class

The typeclass `Show` requires that instances be convertible to `String` (which can then be printed out)

```
class Show a where
  show :: a -> String
```

Indeed, we can test this on different (built-in) types

```
λ> show 2
```

```
"2"
```

```
λ> show 3.14
```

```
"3.14"
```

```
λ> show (1, "two", ([], [], []))
```

```
"(1, \"two\", ([], [], []))"
```

Type Class is a Set of Operations

When we type an expression into ghci, it computes the value and then calls `show` on the result. Thus, if we create a *new* type by

```
data ABC = A | B | C
```

and then create values of the type,

```
λ> let x = A
```

```
λ> :type x
```

```
x :: ABC
```

Type Class is a Set of Operations

but then we cannot view them

```
λ> x
```

```
<interactive>:1:0:
```

```
  No instance for (Show ABC)
```

```
    arising from a use of `print' at <interactive>:1:0
```

```
  Possible fix: add an instance declaration for (Show  
ABC)
```

```
  In a stmt of a 'do' expression: print it
```

Type Class is a Set of Operations

and we cannot compare them!

```
λ> x == x
```

```
<interactive>:1:0:
```

```
  No instance for (Eq ABC)
```

```
    arising from a use of `==' at <interactive>:1:0-5
```

```
  Possible fix: add an instance declaration for (Eq ABC)
```

```
  In the expression: x == x
```

```
  In the definition of `it': it = x == x
```

Again, the previously incomprehensible type error message should make sense to you.

Creating Instances

Tell Haskell how to show or compare values of type **ABC**

By creating instances of **Eq** and **Show** for that type:

instance Eq ABC where

```
(==) A A = True           -- True if both inputs are A
```

```
(==) B B = True           -- ...or B
```

```
(==) C C = True           -- .. or C
```

```
(==) _ _ = False         -- otherwise
```

```
(/=) x y = not (x == y)  -- Test if `x == y` and negate  
result!
```

EXERCISE Lets create an **instance** for **Show ABC**

Automatic Derivation

This is silly: we *should* be able to compare and view `Unshowble` “automatically”!

Haskell lets us *automatically derive* functions for some classes in the standard library.

To do so, we simply dress up the data type definition with

```
data ABC' = A' | B' | C'  
    deriving (Eq, Show) -- tells Haskell to automatically  
generate instances
```

Automatic Derivation

```
data ABC' = A' | B' | C'
```

```
  deriving (Eq, Show) -- tells Haskell to automatically  
generate instances
```

Now we have

```
λ> let x' = A'
```

```
λ> :type x'
```

```
x' :: ABC'
```

```
λ> x'
```

```
A'
```

```
λ> x' == x'
```

```
True
```

```
λ> x' == B'
```

```
False
```


Standard Typeclass Hierarchy

Let us now peruse the definition of the `Num` typeclass.

```
λ> :info Num
```

```
class Num a where
```

```
  (+) :: a -> a -> a
```

```
  (*) :: a -> a -> a
```

```
  (-) :: a -> a -> a
```

```
  negate :: a -> a
```

```
  abs :: a -> a
```

```
  signum :: a -> a
```

```
  fromInteger :: Integer -> a
```

A type `a` is an instance of (i.e. implements) `Num` if there are functions for adding, multiplying, subtracting, negating etc values of that type.

The Ord Typeclass

Another typeclass you've used already is the one for **Ordering** values:

```
λ> :info (<)
class Eq a => Ord a where
  ...
  (<) :: a -> a -> Bool
  ...
```

For example:

```
λ> 2 < 3
True
```

```
λ> "cat" < "dog"
True
```

A type `a` is an instance of (i.e. implements) **Ord** if

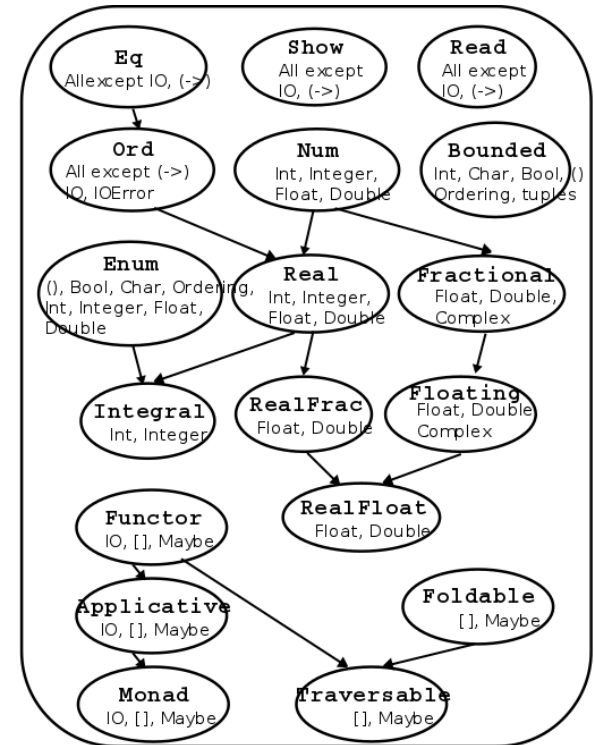
1. It has an instance of **Eq**
2. there are functions for comparing the relative order of values of that type.

Standard Typeclass Hierarchy

In other words in addition to the “arithmetic” operations, we can compare two **Num** values and we can view them (as a **String**.)

Haskell comes equipped with a rich set of built-in classes.

In the picture, there is an edge from **Eq** to **Ord** because for something to be an **Ord** it must also be an **Eq**.



QUIZ

Recall the datatype:

```
data ABC' = A' | B' | C' deriving (Eq, Show)
```

What is the result of:

```
λ> A' < B'
```

- (A) True
- (B) False
- (C) Type error
- (D) Run-time exception



<http://tiny.cc/cse116-ord-ind>

QUIZ

Recall the datatype:

```
data ABC' = A' | B' | C' deriving (Eq, Show)
```

What is the result of:

```
λ> A' < B'
```

- (A) True
- (B) False
- (C) Type error
- (D) Run-time exception



<http://tiny.cc/cse116-ord-grp>

Using Typeclasses

Typeclasses integrate with the rest of Haskell's type system.

Lets build a small library for *Environments* mapping keys *k* to values *v*

```
data Env k v
  = Def v           -- default value `v`
                    -- to be used for "missing" keys
  | Bind k v (Env k v) -- bind key `k` to the value `v`
deriving (Show)
```

An API for Env

Lets write a small API for **Env**

```
-- >>> Let env0 = add "cat" 10.0 (add "dog" 20.0 (Def 0))
```

```
-- >>> get "cat" env0
```

```
-- 10
```

```
-- >>> get "dog" env0
```

```
-- 20
```

```
-- >>> get "horse" env0
```

```
-- 0
```

An API for Env

Ok, lets implement!

-- | 'add key val env' returns a new env that additionally maps `key` to `val`

add :: k -> v -> Env k v -> Env k v

add key val env = ???

-- | 'get key env' returns the value of `key` and the "default" if no value is found

get :: k -> Env k v -> v

get key env = ???

An API for Env

Ok, lets implement!

-- | 'add key val env' returns a new env that additionally maps `key` to `val`

```
add :: k -> v -> Env k v -> Env k v
```

```
add key val env = Bind key val env
```

-- | 'get key env' returns the value of `key` and the "default" if no value is found

```
get :: k -> Env k v -> v
```

```
get key (Def val) = val
```

```
get key (Bind key' val env) | key == key' = val
```

```
get key (Bind key' val env) | otherwise = get key env
```

Constraint Propagation

Lets *delete* the types of `add` and `get` and see what Haskell says their types are!

```
λ> :type get
```

```
get :: (Eq k) => k -> v -> Env k v -> Env k v
```

Haskell tells us that we can use any `k` value as a *key* as long as the value is an instance of the `Eq` typeclass.

How, did GHC figure this out?

- If you look at the code for `get` you'll see that we check if two keys *are equal*!

Exercise

Write an optimized version of

- `add` that ensures the keys are in *increasing* order,
- `get` that gives up and returns the “default” the moment we see a key that's larger than the one we're looking for.

(How) do you need to change the type of `Env`?

(How) do you need to change the types of `get` and `add`?

Explicit Signatures

While Haskell is pretty good about inferring types in general, there are cases when the use of type classes requires explicit annotations (which change the behavior of the code.)

For example, `Read` is a built-in typeclass, where any instance `a` of `Read` has a function

```
read :: (Read a) => String -> a
```

which can parse a string and turn it into an `a`.

That is, `Read` is the *opposite* of `Show`.

QUIZ

What does the expression `read "2"` evaluate to?

- (A) compile time error
- (B) `"2" :: String`
- (C) `2 :: Integer`
- (D) `2.0 :: Double`
- (E) run-time exception



<https://tiny.cc/cse116-read-ind>

QUIZ

What does the expression `read "2"` evaluate to?

- (A) compile time error
- (B) `"2" :: String`
- (C) `2 :: Integer`
- (D) `2.0 :: Double`
- (E) run-time exception



<https://tiny.cc/cse116-read-grp>

Explicit Signatures

Haskell is confused!

- Doesn't know *what type* to convert the string to!
- Doesn't know *which* of the read functions to run!

Did we want an **Int** or a **Double** or maybe something else altogether?

Thus, here an **explicit type annotation** is needed to tell Haskell what to convert the string to:

```
λ> (read "2") :: Int
```

```
2
```

```
λ> (read "2") :: Float
```

```
2.0
```

Note the different results due to the different types.

Creating Typeclasses

Typeclasses are useful for *many* different things.

We will see some of those over the next few lectures.

Lets conclude today's class with a quick example that provides a small taste.

JSON

JavaScript Object Notation or **JSON** is a simple format for transferring data around. Here is an example:

```
{ "name"      : "Elliot Alderson"  
  , "age"      : 28  
  , "likes"    : ["coffee", "hacking"]  
  , "hates"    : [ "e-corp" ]  
  , "lunches" : [ {"day" : "monday",    "loc" : "cafe iveta"}  
                  , {"day" : "tuesday",  "loc" : "cruzn gourmet"}  
                  , {"day" : "wednesday", "loc" : "perk"}  
                  , {"day" : "thursday",  "loc" : "burger."}  
                  , {"day" : "friday",    "loc" : "ray's truck"} ]  
}
```

JSON

In brief, each JSON object is either

- a *base* value like a string, a number or a boolean,
- an (ordered) *array* of objects, or
- a set of *string-object* pairs.

A JSON Datatype

We can represent (a subset of) JSON values with the Haskell datatype

```
data JVal
  = JStr  String
  | JNum  Double
  | JBool Bool
  | JObj  [(String, JVal)]
  | JArr  [JVal]
deriving (Eq, Ord, Show)
```

A JSON Datatype

Thus, the above JSON value would be represented by the `JVal`

```
JObj [("name", JStr "Elliot Alderson")
      ,("age",  JNum 28)
      ,("likes", JArr [ JStr "coffee", JStr "hacking" ])
      ,("hates", JArr [ JStr "e-corp" ])
      ,("lunches", JArr [ JObj [("day", JStr "monday")
                                ,("loc", JStr "cafe iveta")]
                          , JObj [("day", JStr "tuesday")
                                ,("loc", JStr "cruzn gourmet")]
                          , JObj [("day", JStr "wednesday")
                                ,("loc", JStr "perk")]
                          , JObj [("day", JStr "thursday")
                                ,("loc", JStr "burger.")]
                          , JObj [("day", JStr "friday")
                                ,("loc", JStr "ray's truck")]
                          ])
      ])
]
```

Serializing Haskell Values to JSON

Let's write a small library to *serialize* Haskell values as JSON.

We could write a bunch of functions like

```
doubleToJSON :: Double -> JVal  
doubleToJSON = JNum
```

```
stringToJSON :: String -> JVal  
stringToJSON = JStr
```

```
boolToJSON    :: Bool -> JVal  
boolToJSON    = JBool
```

Serializing Collections

But what about collections, namely *lists* of things?

```
doublesToJSON    :: [Double] -> JVal
```

```
doublesToJSON xs = JArr (map doubleToJSON xs)
```

```
boolsToJSON      :: [Bool]   -> JVal
```

```
boolsToJSON xs   = JArr (map boolToJSON xs)
```

```
stringsToJSON    :: [String] -> JVal
```

```
stringsToJSON xs = JArr (map stringToJSON xs)
```

This is **getting rather tedious**

- We are rewriting the same code :(

Serializing Collections (with HOFs)

You could abstract by making the *individual-element-converter* a parameter

```
xsToJSON :: (a -> JVal) -> [a] -> JVal
```

```
xsToJSON f xs = JArr (map f xs)
```

```
xysToJSON :: (a -> JVal) -> [(String, a)] -> JVal
```

```
xysToJSON f kvs = JObject [ (k, f v) | (k, v) <- kvs ]
```

Serializing Collections (with HOFs)

But this is **still rather tedious** as you have to pass in the individual data converter (yuck)

```
λ> doubleToJSON 4
```

```
JNum 4.0
```

```
λ> xsToJSON stringToJSON ["coffee", "hacking"]
```

```
JArr [JStr "coffee", JStr "hacking"]
```

```
λ> xysToJSON stringToJSON [("day", "monday"), ("loc",  
"cafe iveta")]
```

```
JObj [("day", JStr "monday"), ("loc", JStr "cafe iveta")]
```


Serializing Collections (with HOFs)

This gets more hideous when you have richer objects like

```
lunches = [ [{"day", "monday"},      ("loc", "zanzibar")]
            , [{"day", "tuesday"},  ("loc", "farmers market")]
            ]
```

because we have to go through gymnastics like

```
λ> xsToJSON (xysToJSON stringToJSON) lunches
JArr [ JObject [{"day",JStr "monday"}      ,{"loc",JStr "zanzibar"}]
      , JObject [{"day",JStr "tuesday"}    ,{"loc",JStr "farmers market"}]
      ]
```

So much for *readability*

Is it too much to ask for a magical toJSON that *just works*?

Typeclasses To The Rescue

Lets *define* a typeclass that describes types `a` that can be converted to JSON.

```
class JSON a where  
  toJSON :: a -> JVal
```

Now, just make all the above instances of `JSON` like so

```
instance JSON Double where  
  toJSON = JNum
```

```
instance JSON Bool where  
  toJSON = JBool
```

```
instance JSON String where  
  toJSON = JStr
```

Typeclasses To The Rescue

This lets us uniformly write

```
λ> toJSON 4
```

```
JNum 4.0
```

```
λ> toJSON True
```

```
JBool True
```

```
λ> toJSON "hacking"
```

```
JStr "hacking"
```

Bootstrapping Instances

The real fun begins when we get Haskell to automatically bootstrap the above functions to work for lists and key-value lists!

```
instance JSON a => JSON [a] where  
  toJSON xs = JArr [toJSON x | x <- xs]
```

The above says, if `a` is an instance of `JSON`, that is, if you can convert `a` to `JVal` then here's a generic recipe to convert lists of `a` values!

```
λ> toJSON [True, False, True]  
JArr [JBool True, JBool False, JBool True]  
λ> toJSON ["cat", "dog", "Mouse"]  
JArr [JStr "cat", JStr "dog", JStr "Mouse"]
```

or even lists-of-lists!

```
λ> toJSON [["cat", "dog"], ["mouse", "rabbit"]]  
JArr [JArr [JStr "cat", JStr "dog"], JArr [JStr "mouse", JStr "rabbit"]]
```

Bootstrapping Instances

We can pull the same trick with key-value lists

```
instance (JSON a) => JSON [(String, a)] where  
  toJSON kvs = JObject [ (k, toJSON v) | (k, v) <- kvs ]
```

after which, we are all set!

```
λ> toJSON lunches
```

```
JArr [ JObject [ ("day", JStr "monday"), ("loc", JStr "cafe iveta") ]  
      , JObject [ ("day", JStr "tuesday"), ("loc", JStr "cruzn gourmet") ]  
      ]
```

Bootstrapping Instances

It is also useful to bootstrap the serialization for tuples (up to some fixed size) so we can easily write “non-uniform” JSON objects where keys are bound to values with different shapes.

```
instance (JSON a, JSON b) => JSON ((String, a), (String, b)) where  
  toJSON ((k1, v1), (k2, v2)) =  
    JObject [(k1, toJSON v1), (k2, toJSON v2)]
```

```
instance (JSON a, JSON b, JSON c) => JSON ((String, a), (String, b),  
(String, c)) where  
  toJSON ((k1, v1), (k2, v2), (k3, v3)) =  
    JObject [(k1, toJSON v1), (k2, toJSON v2), (k3, toJSON v3)]
```

...

Bootstrapping Instances

Now, we can simply write

```
hs = (("name"    , "Elliot Alderson")
      ,("age"    , 28)
      ,("likes"  , ["coffee", "hacking"])
      ,("hates"  , ["e-corp"])
      ,("lunches", lunches)
      )
```

which is a Haskell value that describes our running JSON example, and can convert it directly like so

```
js2 = toJSON hs
```

That's all folks!
