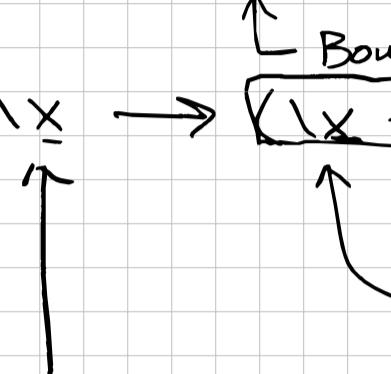


CSE 114A Lecture 4

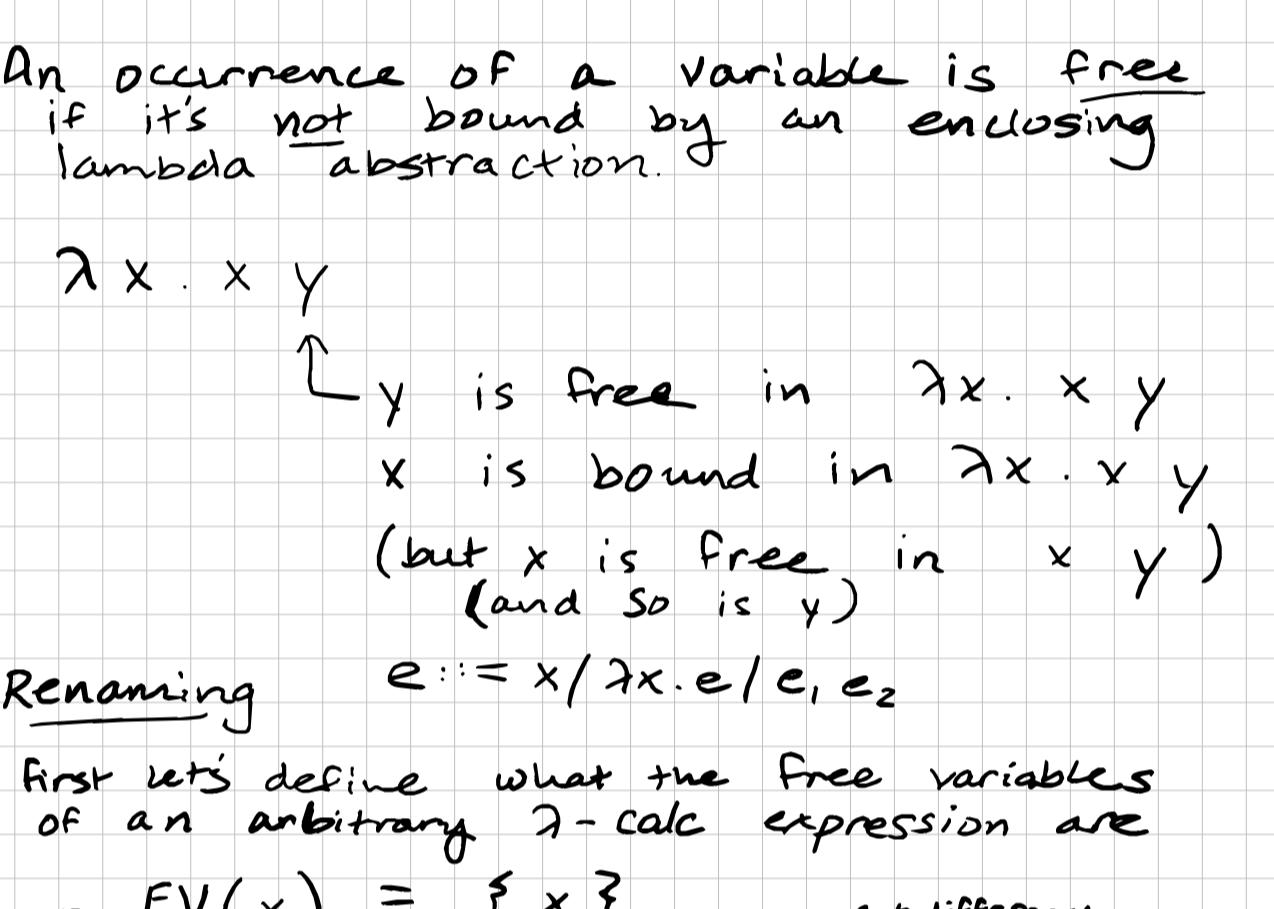
- ✓ - announcement
- wrap-up of lambda calc:
 - ✓ - Scope of a variable; renaming
 - ✓ - Normal form; nontermination (omega)
 - pairs (structs, tuples, ...)
 - Why is ADD defined how it is?
 - Recursion!
- zooming out: the big picture.
Why discuss λ -calc?

Scope and renaming.

Scope of a variable: the part of a program where the variable is visible.



The body of this function, e , is the scope of x . Any occurrence of x in e is bound by the binder λx .



An occurrence of a variable is free if it's not bound by an enclosing lambda abstraction.

$$\lambda x. x y$$

y is free in $\lambda x. x y$
 x is bound in $\lambda x. x y$
(but x is free in $x y$)
(and so is y)

Renaming $e ::= x / \lambda x. e / e_1 e_2$

First let's define what the free variables of an arbitrary λ -calc expression are

$$\begin{aligned} - FV(x) &= \{x\} && \text{set difference} \\ - FV(\lambda x. e) &= FV(e) - \{x\} \\ - FV(e_1 e_2) &= FV(e_1) \cup FV(e_2) \end{aligned}$$

This gives us a policy for how to do renaming!

$$\lambda x. e \xrightarrow{\alpha} \lambda y. e[x := y]$$

where y is not in $FV(e)$

$$\lambda x. x = \lambda y. y = \lambda z. z, \text{ for example.}$$

In Elsa, much like $=b>$ is beta reduction, (aka substitution), $=a>$ is alpha reduction (aka renaming).

$$\lambda x. \lambda y. x \xrightarrow{\alpha} \lambda x. \lambda y. (\lambda x. x)$$

you could do this, but it's easy to get confused.

$$(\lambda x. \lambda y. (\lambda q. q)) \ x \xrightarrow{\beta} \lambda y. (\lambda q. q)$$

Normal Form

A "reducible expression" (or redex for short) is an expression of the form $(\lambda x. e_1) e_2$ (in other words, the left side of the substitution rule!).

If you have any redexes anywhere in a λ -calc expression, it's not in normal form.

If you have no redexes, it's in normal form.

An expression in normal form is a value (and can't be further reduced).

$$\lambda x. y z$$

an application of y to z

$$(\lambda y. z \rightarrow (\lambda x \rightarrow y z)) \ \text{bulbasaur man}$$

$=b>$

$$(\lambda z \rightarrow (\lambda x \rightarrow \text{bulbasaur } z)) \ \text{man}$$

$=b>$

$$\lambda x \rightarrow \underbrace{\text{bulbasaur man}}$$

↑
This might eventually be a redex, but right now it can't be reduced.

Do all λ -calc expressions have a normal form?
(Do they all evaluate to a value?)

No!

$$(\lambda x. \boxed{x \quad x}) (\lambda x. x \quad x)$$

$x \xrightarrow{\text{alpha red}} x$

The omega
combinator.

$$\xrightarrow{B} (\lambda x. x \quad x) (\lambda x. x \quad x)$$

oh no!

We'll never get to a value.