# CSE 116: Fall 2019 Introduction to Functional Programming

#### Intro to Haskell

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#### What is Haskell?

- A typed, lazy, purely functional programming language
  - Haskell = λ-calculus +
    - Better syntax
    - Types
    - Built-in features
      - Booleans, numbers, characters
      - Records (tuples)
      - Lists
      - Recursion

- ...

# Why Haskell?

- Haskell programs tend to be simple and correct
- Quicksort in Haskell

```
sort [] = []
sort (x:xs) = sort ls ++ [x] ++ sort rs
    where
    ls = [ l | l <- xs, l <= x ]
    rs = [ r | r <- xs, x < r ]</pre>
```

- Goals for this week
  - Understand the above code
  - Understand what typed, lazy, and purely functional means (and why you care)

# Haskell vs λ-calculus: Programs

- A program is an expression (not a sequence of statements)
- It evaluates to a value (it does not perform actions)
  - λ: (\x -> x) apple -- =~> apple
  - Haskell: (\x -> x) "apple" -- =~> "apple"

#### Haskell vs λ-calculus: Functions

- Functions are first-class values:
  - can be *passed as arguments* to other functions
  - can be *returned as results* from other functions
  - can be partially applied (arguments passed one at a time)

```
(\x -> (\y -> x (x y))) (\z -> z + 1) 0 -- =\sim> 2
```

• BUT: unlike  $\lambda$ -calculus, not everything is a function!

## Haskell vs λ-calculus: top-level bindings

- Like in Elsa, we can name terms to use them later
- Elsa:

```
let T = \langle x y - \rangle x
let F = \langle x y \rightarrow y \rangle
let PAIR = \xy -> \b -> \xy
let FST = \protect\ p \rightarrow p T
let SND = \protect\ p \rightarrow p F
eval fst:
 FST (PAIR apple orange)
 =~> apple
```

## Haskell vs λ-calculus: top-level bindings

- Like in Elsa, we can name terms to use them later
- Haskell:

```
haskellIsAwesome = True
pair = \x y -> \b -> if b then x else y
fst = \p -> p haskellIsAwesome
snd = \p -> p False

-- In GHCi:
> fst (pair "apple" "orange") -- "apple"
```

- The names are called top-level variables
- Their definitions are called top-level bindings

You can define function bindings using equations:

```
pair x y b = if b then x else y -- pair = \xspace x y b -> \dots
fst p = p True -- fst = \xspace p -> \dots
snd p = p False -- snd = \xspace p -> \dots
```

 A single function binding can have multiple equations with different patterns of parameters:

- The first equation whose pattern matches the actual arguments is chosen
- For now, a pattern is:
  - a variable (matches any value)
  - or a value (matches only that value)

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Same as:

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Same as:

# QUIZ: Pair

Which of the following definitions of pair is incorrect? \*

```
A. pair x y = b \rightarrow \mathbf{if} b \mathbf{then} x \mathbf{else} y
B. pair x = y b \rightarrow if b then x else y
C.
pair x _ True = x
pair _ y _ = y
D.
pair x y b = x
pair x y False = y
```



E. all of the above

http://tiny.cc/cse116-pair-ind

# QUIZ: Pair

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### Equations with guards

An equation can have multiple guards (Boolean expressions):

Same as:

#### Recursion

• Recursion is built-in, so you can write:

• Or you can write:

```
sum 0 = 0
sum n = n + sum (n - 1)
```

#### Scope of variables

Top-level variables have global scope

Or you can write:

```
-- What does f compute?
f 0 = True
f n = g (n - 1) -- mutual recursion!
g 0 = False
g n = f (n - 1) -- mutual recursion!
```

• Answer: f is is Even, g is is 0dd

### Scope of variables

Is this allowed?

```
haskellIsAwesome = True
haskellIsAwesome = False -- changed my mind
```

 Answer: no, a variable can be defined once per scope; no mutation!

#### Local variables

 You can introduce a new (local) scope using a 1etexpression

Syntactic sugar for nested 1et-expressions:

#### Local variables

 If you need a variable whose scope is an equation, use the where clause instead:

What would Elsa say?

```
let FNORD = ONE ZERO
```

- **Answer**: Nothing. When evaluated, it will crunch to something, but it will be nonsensical.
  - $\lambda$ -calculus is **untyped**.

What would Python say?

```
def fnord():
   return 0(1)
```

- Answer: Nothing. When evaluated will cause a runtime error.
  - Python is dynamically typed

• What would Java say?

```
void fnord() {
  int zero;
  zero(1);
}
```

- Answer: Java compiler will reject this.
  - Java is **statically typed**.

- In Haskell every expression either has a type or is illtyped and rejected statically (at compile-time, before execution starts)
  - like in Java
  - unlike λ-calculus or Python

```
fnord = 1 0 -- rejected by GHC
```

#### Type Annotations

• You can annotate your bindings with their types using ::, like so:

#### Type Annotations

```
-- | This is a word-size integer
rating :: Int
rating = if haskellIsAwesome then 10 else 0

-- | This is an arbitrary precision integer
bigNumber :: Integer
bigNumber = factorial 100
```

- If you omit annotations, GHC will infer them for you
  - Inspect types in GHCi using :t
  - You should annotate all top-level bindings anyway!
     (Why?)

#### **Function Types**

- Functions have arrow types
  - $\x$  -> e has type A -> B
  - If e has type B, assuming x has type A
- For example:

```
> :t (\x -> if x then 'a' else 'b')
(\x -> if x then 'a' else 'b') :: Bool -> Char
```

#### **Function Types**

You should annotate your function bindings:

```
sum :: Int -> Int

sum 0 = 0

sum n = n + sum (n - 1)
```

With multiple arguments:

```
pair :: String -> (String -> (Bool -> String))
pair x y b = if b then x else y
```

Same as:

```
pair :: String -> String -> Bool -> String
pair x y b = if b then x else y
```

# QUIZ: Type of Pair

With pair :: String -> String -> Bool -> String, what would GHCi say

```
>:t pair "apple" "orange"
```

- **A.** Syntax error
- **B.** The term is ill-typed
- C. String
- D. Bool -> String
- E. String -> String -> Bool -> String



http://tiny.cc/cse116-tpair-ind

# QUIZ: Type of Pair

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>:t pair "apple" "orange"
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http://tiny.cc/cse116-tpair-grp

#### Lists

- A list is
  - either an empty list

```
[] -- pronounced "nil"
```

- or a *head element* attached to a *tail list* 

```
x:xs -- pronounced "x cons xs"
```

#### Terminology: constructors and values

```
-- A list with zero elements
                  -- A list with one element: 1
1:[]
(:) 1 []
                  -- Same thing: for any infix op,
                  -- (op) is a regular function!
1:(2:(3:(4:[]))) -- A list with four elements: 1, 2, 3, 4
                  -- Same thing (: is right associative)
1:2:3:4:[]
[1,2,3,4]
                  -- Same thing (syntactic sugar)
```

#### Lists

- [] and (:) are called the list constructors
- We've seen constructors before:
  - True and False are Bool constructors
  - 0, 1, 2 are... well, it's complicated, but you can think of them as Int constructors
  - these constructions didn't take any parameters, so we just called them *values*
- In general, a **value** is a constructor applied to *other* values (e.g., *list values* on previous slide)

## Type of a list

- A list has type [A] if each one of its elements has type A
- Examples:

## Functions on lists: range

There is also syntactic sugar for this!

```
[1..7] -- [1,2,3,4,5,6,7]
[1,3..7] -- [1,3,5,7]
```

## Functions on lists: length

```
-- | Length of the List
length :: ???
length xs = ???
```

### Pattern matching on lists

```
-- | Length of the List
length :: [Int] -> Int
length [] = 0
length (_:xs) = 1 + length xs
```

- A pattern is either a variable (incl. \_) or a value
- A pattern is
  - either a variable (incl. \_)
  - or a *constructor* applied to other *patterns*
- Pattern matching attempts to match values against patterns and, if desired, bind variables to successful matches.

### **QUIZ:** Patterns

Which of the following is not a pattern? \*

- A. (1 : xs)
- O B. (\_:\_:\_)
- O. [x]
- O. [1+2, x, y]
- E. all of the above



http://tiny.cc/cse116-pattern-ind

## QUIZ: Patterns (wrong url)

Which of the following is not a pattern? \*

- A. (1 : xs)
- O B. (\_:\_:\_)
- O. [x]
- O. [1+2, x, y]
- E. all of the above



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### Some useful library functions

```
-- | Is the list empty?
null :: [t] -> Bool
-- | Head of the list
head :: [t] -> t -- careful: partial function!
-- | Tail of the list
tail :: [t] -> [t] -- careful: partial function!
-- | Length of the list
length :: [t] -> Int
-- | Append two lists
(++) :: [t] -> [t] -> [t]
-- | Are two lists equal?
(==) :: [t] -> [t] -> Bool
```

You can search for library functions (by type!) at hoogle.haskell.org

#### **Pairs**

isEmpty p

```
myPair :: (String, Int) -- pair of String and Int
   myPair = ("apple", 3)
• (,) is the pair constructor
   -- Field access using library functions:
   whichFruit = fst myPair -- "apple"
   howMany = snd myPair -- 3
   -- Field access using pattern matching:
   is Empty (x, y) = y == 0
                                               You can use pattern
                                                matching not only
   -- same as:
                                                in equations, but
                     = \setminus (x, y) \rightarrow y == 0
   isEmpty
                                                also in \lambda-bindings
                                                and let-bindings!
   -- same as:
```

= let (x, y) = p in y == 0

### Pattern matching with pairs

 Is this pattern matching correct? What does this function do?

### Pattern matching with pairs

 Is this pattern matching correct? What does this function do?

 Answer: a list of pairs represents key-value pairs in a dictionary; f performs lookup by key

### **Tuples**

- Can we implement triples like in  $\lambda$ -calculus?
- Sure! But Haskell has native support for n-tuples:

```
myPair :: (String, Int)
myPair = ("apple", 3)
myTriple :: (Bool, Int, [Int])
myTriple = (True, 1, [1,2,3])
my4tuple :: (Float, Float, Float)
my4tuple = (pi, sin pi, cos pi, sqrt 2)
-- And also:
myUnit :: ()
myUnit = ()
```

### List comprehensions

A convenient way to construct lists from other lists:

```
[toUpper c | c <- s] -- Convert string s to upper case
```

# Quicksort in Haskell

```
sort [] = []
sort (x:xs) = sort ls ++ [x] ++ sort rs
    where
    ls = [ l | l <- xs, l <= x ]
    rs = [ r | r <- xs, x > r ]
```

### What is Haskell?

• A typed, lazy, purely functional programming language

### Haskell is statically typed

- Every expression either has a type, or is ill-typed and rejected at compile time
- Why is this good?
  - catches errors early
  - types are contracts (you don't have to handle illtyped inputs!)
  - enables compiler optimizations

### Haskell is purely functional

- **Functional** = functions are *first-class values*
- Pure = a program is an expression that evaluates to a value
  - No side effects! unlike in Python, Java, etc:

 in Haskell, a function of type Int -> Int computes a single integer output from a single integer input and does nothing else

### Haskell is purely functional

- Referential transparency: The same expression always evaluates to the same value
  - More precisely: In a scope where x1, ..., xn are defined, all occurrences of e with
     FV(e) = {x1, ..., xn} have the same value
- Why is this good?
  - easier to reason about (remember x++ vs ++x in C?)
  - enables compiler optimizations
  - especially great for parallelization (e1 + e2: we can always compute e1 and e2 in parallel!)

### Haskell is lazy

- An expression is evaluated only when its result is needed
- Example: take 2 [1 .. (factorial 100)]

```
take 2 ( upto 1 (factorial 100))

=> take 2 ( upto 1 933262154439...)

=> take 2 (1:(upto 2 933262154439...)) -- def upto

=> 1: (take 1 ( upto 2 933262154439...)) -- def take 3

=> 1: (take 1 (2:(upto 3 933262154439...)) -- def upto

=> 1:2:(take 0 ( upto 3 933262154439...)) -- def take 3

=> 1:2:[] -- def take 1
```

\_

### Haskell is lazy

- Why is this good?
  - Can implement cool stuff like infinite lists: [1..]

- encourages simple, general solutions
- but has its problems too :(

### That's all folks!