## CSE 116, Fall 2019 Midterm

| Section | Points | Score |
| :---: | :---: | :---: |
| Part I | 40 points |  |
| Part II | 56 points |  |
| Total | 96 points |  |

## Instructions

- You have 95 minutes to complete this exam.
- This exam is closed book. You may use one double-sided page of notes, but no other materials.
- Avoid seeing anyone else's work or allowing yours to be seen.
- Do not communicate with anyone but an exam proctor.
- To ensure fairness (and the appearance thereof), proctors will not answer questions about the content of the exam. If you are unsure of how to interpret a problem description, state your interpretation clearly and concisely. Reasonable interpretations will be taken into account by graders.

NAME: $\qquad$

CruzID: $\qquad$ @ucsc.edu

## Part I: Lambda calculus

1. [16pts] Use $\beta$-reductions to evaluate the following lambda term to a normal form.

(B) $(\backslash x$ y $\rightarrow(y$ x $)(\backslash p q->p))(\backslash i->i)(\backslash j->j)$
2. [8pts] For each bound occurrence of a variable in the following lambda terms, draw an arrow pointing to its binder. For each free occurrence, draw a circle around the variable.
(A) ( $\begin{aligned} & a \\ & -> \\ & b\end{aligned}\left(\begin{array}{llll}\backslash b & a & -> & b\end{array}\right)$
$(B)(\backslash p q \quad r \rightarrow(p(\backslash q \quad p \rightarrow(r q)))(q \quad p))$
3. [16pts] Fill in a lambda calculus expression for each blank in the program below to define a function PROD where (PROD $n$ ) returns the product of numbers between $n$ and one. You may use any of the functions defined on the Lambda Calculus Cheat Sheet on the back page. Any other helper functions you must define yourself. Your implementation may assume that PROD is never called with ZERO.
```
let PROD1 = \f n -> ITE _____(A)
    _______
        (B)
    (C)
```

let $P R O D=$ $\qquad$ (D) $\qquad$
(A)
(C)
(D)

## Part II: Haskell

4. [5pts] What does the following Haskell program evaluate to?
```
let f = (\x -> \y -> x + y)
    g = f 5
    h = \f n -> f (f n)
in
    h g 3
```

(a) Type Error
(b) 8
(c) 13
(d) $\backslash f \rightarrow f(f 8)$
(e) 16
5. [5pts] What is the most general type of the Haskell function foo?
foo $\operatorname{bar}(x, y)$
| bar $x=y++y$
| otherwise = y
(a) (a -> b) -> (a,b) -> [b]
(b) (String -> Bool) -> (String, String) -> String
(c) (a -> Bool) -> (a, a) -> [a]
(d) (Bool -> a) -> [b] -> [b]
(e) (a -> Bool) -> (a, [b]) -> [b]

For the following questions, consider the data types defined below.

```
data TrickOrTreat = Trick Prank | Treat Candy
data Prank = Prank { desc :: String, legal :: YNM }
data YNM = Yes | No | Maybe
data Candy = Candy { pieces :: Int, kind :: Kind, rating :: Int}
data Kind = Chocolate | HardCandy | Gummies
```

6. [21pts] A case expression is exhaustive if all possible values are matched by at least one pattern. A pattern is overlapping if previous patterns match all values it matches. Assume $t$ has type TrickOrTreat and do the following:

- For each pattern in each case, provide a value that matches on the pattern.
- If the pattern is overlapped by previous patterns, write "N/A".
- Determine if the case expression is exhaustive and circle Exhaustive or Non-exhaustive as appropriate.
- For non-exhaustive case expressions, write a value that does not match any of its patterns.
(a)

```
        case t of
    Trick (Prank _ Yes) -> ()
    Trick (Prank d No) -> ()
```

    ---------------------------
    -------------------------
    -------------------------
    [ Exhaustive / Non-exhaustive ]
(b) case $t$ of


(c)
$\square$
$\qquad$



(d)
case $t$ of

Treat c -> ()

Treat (Candy $n$ Chocolate r) -> ()

Trick p -> ()
[ Exhaustive / Non-exhaustive ]

Treat (Candy _ k r) | r > $3->$ ()

Trick (Prank d l) -> ()

Treat c | (rating c) < 3 -> ()
[ Exhaustive / Non-exhaustive ]

Trick (Prank "snakes on plane" No) -> ()

Treat (Candy 99 Gummies r) -> ()
[ Exhaustive / Non-exhaustive ]
7. Consider a binary search tree where each internal node has a value $i$ and two child subtrees. The left child contains all nodes with values less than or equal to $i$, and the right child contains all nodes with values strictly greater than $i$.


Figure 1: A binary search tree
(A) [5pts] Using the following ADT, create a Tree that represents the binary search tree in Figure 1.

```
data Tree = Leaf | Node Int Tree Tree
```

(B) $[10 \mathrm{pts}]$ Define the function max, which returns the highest value in a binary search tree represented by a Tree value, or 0 if the tree is empty. For example, if $t$ is the tree in Figure 1, then max $t$ evaluates to 14. Full credit requires an efficient traversal of the tree (e.g., not an exhaustive one). You may define helper functions if desired, but you may not use any library functions except $==,>=$, or $<=$.

```
max :: Tree -> Int
```

(C) $[10 \mathrm{pts}]$ Define the function contains, which returns True if a number n is contained in a binary search tree $t$. For example, if $t$ is the tree in Figure 1, then contains 7 t returns True, but contains 5 t returns False. Full credit requires an efficient traversal of the tree (e.g., not an exhaustive one). You may define helper functions if desired, but you may not use any library functions except $==,>=$, or $<=$. contains :: Int -> Tree -> Bool

## 1 Lambda calculus cheat sheet

```
-- Booleans
let TRUE =\x y -> x
let FALSE = \x y -> y
let ITE = \b x y -> b x y
let NOT = \b x y -> b y x
let AND = \b1 b2 -> ITE b1 b2 FALSE
let OR = \b1 b2 -> ITE b1 TRUE b2
-- Numbers
let ZERO = \f x-> x
let ONE = \f x -> f x
let TWO = \f x -> f (f x)
let THREE = \f x -> f (f (f x))
let FOUR = \f x -> f (f (f (f x))
let FIVE = \f x -> f (f (f (f (f x))
-- Arithmetic
let INC = \n f x m f (nf x )
let ADD = \n m -> n INC m
let MUL = \n m -> n (ADD m) ZERO
let ISZ = \n -> -- return TRUE if n == 0 --
let DECR = \n -> -- decrement n by one --
let EQL = \a b -> -- return TRUE if a == b, otherwise FALSE --
-- Recursion ------------------------------------
let FIX = \stp -> (\x -> stp ( 
```

